Week 2: Lab

Collaboration level 0 (no restrictions). Open notes.

- 1. Prove or disprove: f = O(g) implies that g = O(f). What we expect: If true, a justification that holds for any functions f, g and for all $n > n_0$. If false, a counterexample.
- 2. Prove or disprove: f = O(g) implies that $g = \Omega(f)$. What we expect: If true, a justification that holds for any functions f, g and for all $n > n_0$. If false, a counterexample.
- 3. Find the order of growth of the following functions. For each function, give its $\Theta()$.
 - (a) $n \lg \lg n + n \lg n + \sqrt{n} \lg^2 n$
 - (b) $\sqrt{n} \lg n + n$
 - (c) $n^2 + \sqrt{n} \lg^3 n$
 - (d) $3^{\lg n} + n^2 + n \lg n$
 - (e) $\sqrt{3}^{\lg n} + n^2 + n \lg n$
 - (f) $2^n + 2^{2n}$
 - (g) $2^{\lg n} + \lg n^2$
 - (h) $(\lg n)^{\lg n} + n^3$
- 4. Arrange the following functions in ascending order of growth rate. For each pair of consecutive functions, give a brief justification on why they are in this order (For e.g., if you ordered A, B, C, you need to justify that A = O(B) and B = O(C)).

$$2^{\sqrt{\log n}}, 2^n, n^{4/3}, n(\log n)^3, n^{\log n}, 2^{2^n}, 2^{n^2}$$

- 5. Express the worst case running time of Insertion Sort as a sum and find its asymptotic order of growth (using the formula discussed this week).
- 6. Consider the **element uniqueness problem**: check whether all the elements in a given array are distinct. The problem can be solved by the following straightforward algorithm:

1

```
UniqueElements (A[0..n-1]

1  // Checks whether all the elements in a given array are distinct

2  // Output: Returns "true" if all elements are distinct, and "false" otherwise

3  For i = 0 to n - 2

4  For j = i + 1 to n - 1

5  A[i] == B[j] return false

6  return true
```

What is the running time (best-case, worst-case) of the algorithm, as a function of the n? (hint: express as a sum and use the formula discussed).

7. Given two *n*-by-*n* matrices A and B, by definition, their product C = AB is an *n*-by-*n* matrix C where

$$C[i,j] = A[i,0]B[0,j] + A[i,1]B[1,j] + \ldots + A[i,k]B[k,j] + \ldots + A[i,n-1]B[n-1,j]$$

The straightforward algorithm to compute the product is given below:

```
MATRIXMULTIPLICATION(A[0..n-1,0..n-1], B[0..n-1,0..n-1])

1  // Multiplies two square matrices of order n by the definition-based algorithm

2  // Output: Matrix C = AB

3  For i = 0 to n - 1

4  For j = 0 to n - 1

5  // computes C[i,j]

6  C[i,j] = 0

7  For k = 0 to n - 1

8  C[i,j] + a[i,k] \times B[k,j]
```

What is the running time (best-case, worst-case) of the algorithm, as a function of the order of the matrix, n?

8. The following algorithm finds the number of binary digits in the binary representation of a positive decimal integer.

¹If you haven't seen matrices before, you can assume a matrix is a 2-dimensional array.

```
BINARY(n)

1  // Input: A positive decimal integer n

2  // Output: The numbr of binary digits in n's binary representation

3  count = 1

4  While n > 1

5   count++

6   n = n/2

7  return count
```

What is the running time (best-case, worst-case) of the algorithm, as a function of n?

9. Consider the following algorithm:

```
Mystery(n)

1  // Input: A positive decimal integer n

2  s = 0

3  For i = 1 to n

4  s = s + i \times i

5  return s
```

- (a) What does this algorithm compute?
- (b) What is the running time of the algorithm, as a function of n?
- 10. Assume that we have an algorithm whose running time is f(n) microseconds. Determine the largest size of a problem that can be solved by the algorithm in: (a) 1 second; (b) 1 hour; (c) 1 month.

(a)
$$f(n) = \lg n$$
 (b) $f(n) = n^{10}$ (c) $f(n) = 2^n$.

Optional

11. You are presented with 9 marbles. All of the marbles look identical i.e. same shape, color, and dimensions(except for weight). However, 8 of the 9 marbles have exactly the same weight; the last marble is heavier. The only tool you have to measure weights is an old fashioned balance scale. You are only allowed to use the scale 2 times. How do you find the one marble that is not the same weight as the others?

Generalize to n marbles: find it using the scale $\log_3 n$ times. (note: interview question)

12. You are given a set of n points on a circle in the plane. Come up with an algorithm that determines if there exists a pair of points that are antipodal (two points are antipodal if they are diametrically opposite). Analyze its running time (note: interview question).